

The Subtree Polynomial

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Joint work with Jason Brown (Dalhousie University)

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COMPLEX SUBTREE ROOTS

REAL SUBTREE ROOTS

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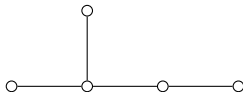
- ▶ Alternatively,

$$\Phi_T(x) = \sum_{k=1}^n s_k(T) x^k,$$

where $s_k(T)$ is the number of subtrees of T of order k .

AN EXAMPLE

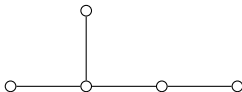
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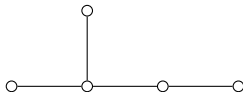
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$$\Phi_T(x) = 5x$$

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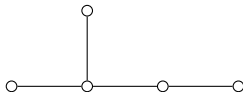
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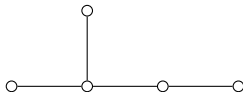
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$$\Phi_T(x) = 5x + 4x^2 + 4x^3$$

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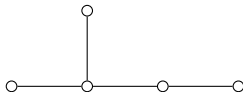
Consider the following tree T :



$$\Phi_T(x) = 5x + 4x^2 + 4x^3 + 3x^4$$

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$$M_T = \frac{\Phi_T'(1)}{\Phi_T(1)}$$

is the *mean subtree order* of T .

- ▶ First studied by Jamison.
- ▶ Less obvious: $-\Phi_T(-1)$ is the *independence number* of T (Jamison, 1987).

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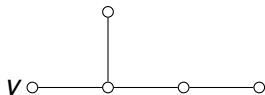
$$\Phi_{T,v}(x) = \sum_{S \in \mathcal{S}_v} x^{|V(S)|}.$$

Evidently, we have:

$$\Phi_T(x) = \Phi_{T,v}(x) + \Phi_{T-v}(x).$$

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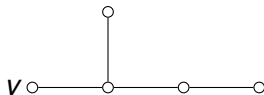
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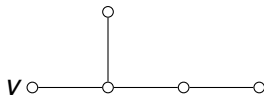
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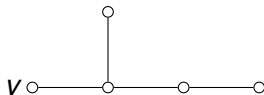
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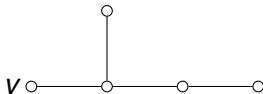
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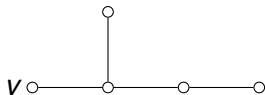
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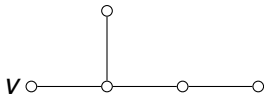
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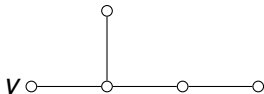


$$\Phi_{T,v}(x) = x + x^2 + 2x^3 + 2x^4 + x^5$$

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Notice: Many of the “large” subtrees of T contain v , while many of the “small” subtrees of T do not contain v .

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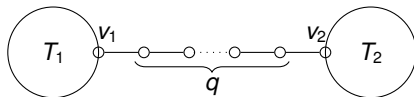
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- ▶ We give a short proof of this result.

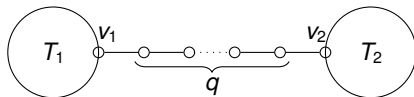
CSIKVÁRI'S GENERALIZED TREE SHIFT

Let T be a tree, and let v_1 and v_2 be non-leaf vertices of T connected by a path whose internal vertices all have degree 2.

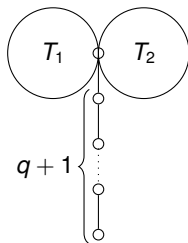


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The following tree T' is said to be obtained from T by a *generalized tree shift*:



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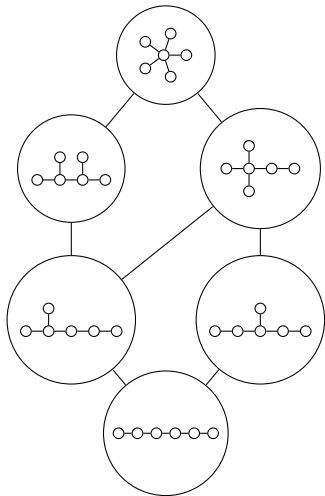
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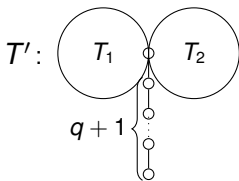
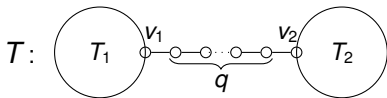
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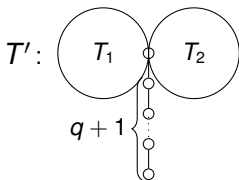
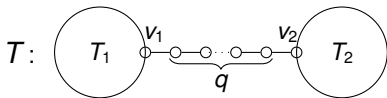


Theorem: Let T be a tree of order $n \geq 4$, and let T' be a tree obtained from T by a generalized tree shift. For every $k \in \{3, \dots, n-1\}$, $s_k(T) < s_k(T')$.

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Sketch of Proof: Show that

$$\Phi_{T'}(x) - \Phi_T(x) = \frac{1}{x} \cdot \left[\sum_{i=0}^q x^i \right] \cdot [\Phi_{T_1, v_1}(x) - x] \cdot [\Phi_{T_2, v_2}(x) - x].$$

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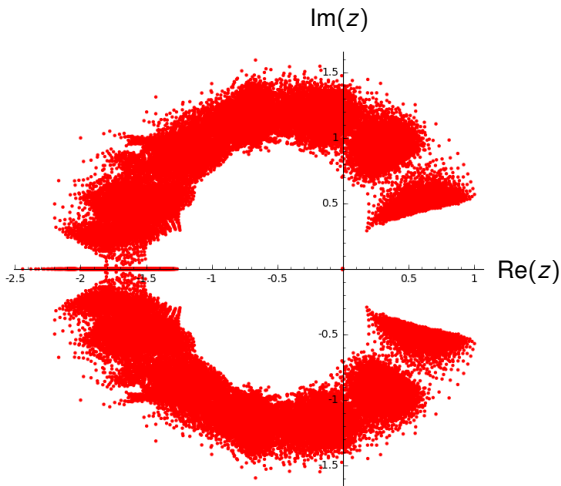
OPEN PROBLEMS

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The subtree roots of all trees of order at most 14.

BOUNDING THE MODULUS

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- ▶ This bound is tight! The only tree with a subtree root of modulus $1 + \sqrt[3]{3}$ is the star $K_{1,3}$.

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This is proven using a rather technical strong induction. The key is a good lower bound on $|\Phi_{T,v}(z)|$.

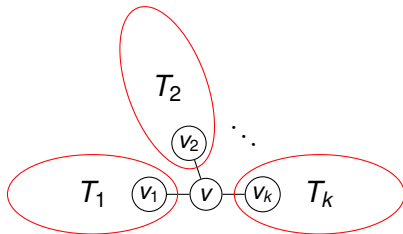
Lemma: Let T be a tree of order n with vertex v . If $|z| \geq 2$, then

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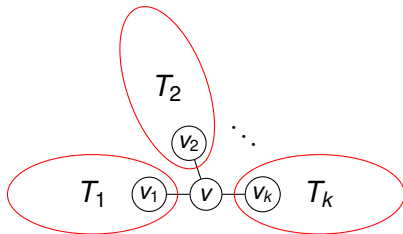
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Key ingredient of proof:



$$\Phi_{T,v}(z) = z \cdot \prod_{i=1}^k [1 + \Phi_{T_i, v_i}(z)]$$

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- ▶ The subtree polynomial has no complex roots of modulus greater than $1 + \sqrt[3]{3}$
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- ▶ Question: What about the interval $(-1 - \sqrt[3]{3}, 0)$?

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- ▶ **Useful identity:** $x\Phi'_T(x) = \sum_{v \in V(T)} \Phi_{T,v}(x)$
- ▶ **Claim:** $\forall v \in V(T), x \in (-1, 0) \Rightarrow \Phi_{T,v}(x) \in (-1, 0)$.
- ▶ Claim is proved using induction and the nice recursive formula for $\Phi_{T,v}(x)$:

$$\Phi_{T,v}(x) = x \cdot \prod_{i=1}^k [1 + \Phi_{T_i, v_i}(x)]$$

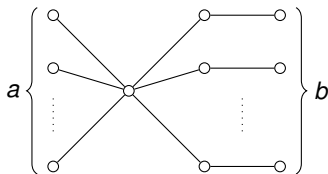
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Theorem: The closure of the collection of all real subtree roots contains the interval $[-2, -1]$.

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Proof technique: We show that the collection of real subtree roots of the following family of trees is dense in $[-2, -1]$:



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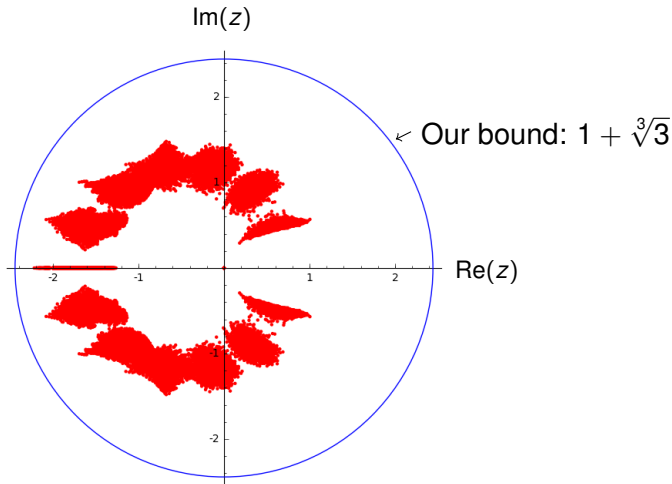
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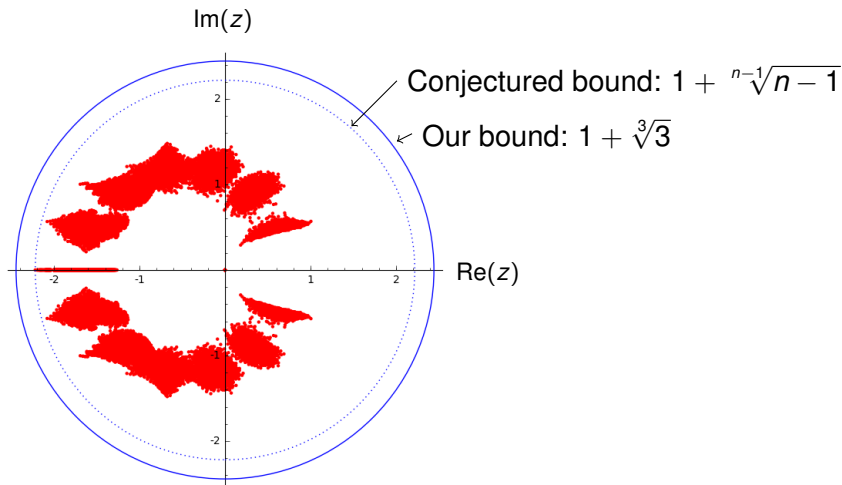
OPEN PROBLEMS

Find a tighter bound on the subtree roots of all trees of order n .



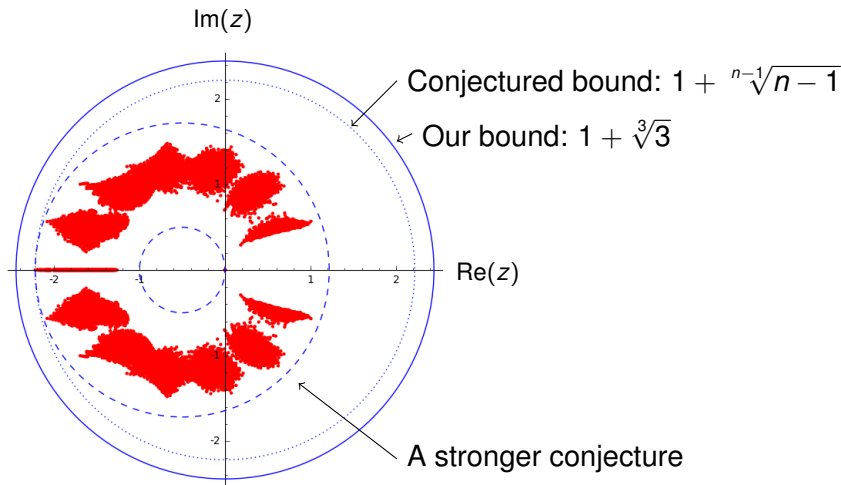
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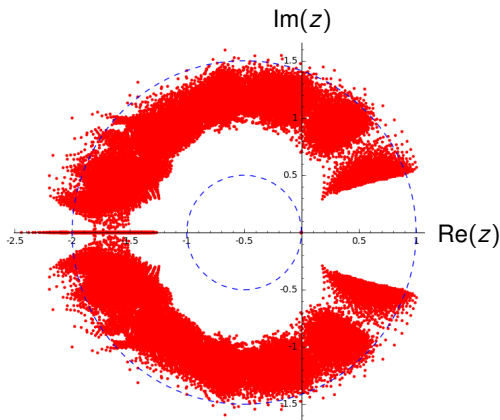
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CLOSURE

Does the closure of the collection of subtree roots contain the entire annulus $\frac{1}{2} \leq |z + \frac{1}{2}| \leq \frac{3}{2}$?



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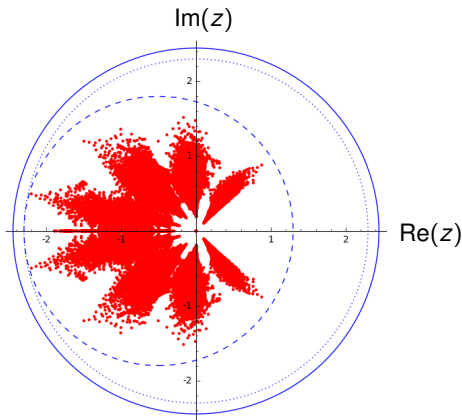
- ▶ Consider “subtrees” of *graphs*.

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- ▶ For which trees is the sequence of coefficients $s_2(T), s_3(T), \dots, s_n(T)$ unimodal?
- ▶ In particular, is this sequence unimodal when T has no vertices of degree 2?
 - ▶ Fact: This sequence is NOT log-concave for all such trees.

INTRODUCTION
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COUNTING SUBTREES
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COMPLEX SUBTREE ROOTS
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REAL SUBTREE ROOTS
○○○○○

OPEN PROBLEMS
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Thank you!